

Nonlinear Systems, Theory III

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Now let the x_k 's represent inputs at k points, spaced by Δt .

$$R(t) = a_0 + \sum_k q_k s(t - k\Delta t) + \sum_k b_k (s(t - k\Delta t)^2 - P) \\ + \sum_{k < l} c_{kl} s(t - k\Delta t) s(t - l\Delta t) \\ \dots$$

where $P = \langle s(t)^2 \rangle$.

The P -term in $b_k \dots$ comes from the fact that $\langle s(t - k\Delta t)^2 \rangle \neq 0$, so the b -term would otherwise contribute a constant, confounding the a_0 -term.

Does this make sense as $\Delta t \rightarrow 0$?

$s(t)$ gets sampled more & more densely. But expect that the influence of $s(t - \tau)$ on $R(t)$ depends on τ , not the lag number. i.e.,

expect $a_k \rightarrow \Delta t A(k\Delta t)$. Think of $s(t) = \begin{array}{|c|c|c|c|}\hline & & & \\ \hline \end{array}$, resample each Δt .

Similarly, there are $(\frac{1}{\Delta t})^2$ as many samples for c_{kl} , so expect that

$$c_{kl} \rightarrow (\Delta t)^2 C(k\Delta t, l\Delta t).$$

$$b_k \rightarrow (\Delta t)^2 b(k\Delta t)$$

\uparrow
for dimensional correctness.

(3)

This motivates: $K_0 = \alpha_0$

$$K_1(\gamma) = A(\gamma)$$

$$K_2(\gamma_1, \gamma_2) = \begin{cases} \frac{1}{2} C(\gamma_1, \gamma_2) & \gamma_1 < \gamma_2 \\ \frac{1}{2} C(\gamma_2, \gamma_1) & \gamma_1 > \gamma_2 \\ b(\gamma_1) & \gamma_1 = \gamma_2 \end{cases}$$

$$R(t) = K_0 + \sum_k K_1(k\Delta t) s(t - k\Delta t) \cdot \Delta t$$

$$+ \sum_k K_2(k\Delta t, k\Delta t) (s(t - k\Delta t)^2 - P) (\Delta t)^2$$

$$+ \sum_{k \neq l} K_2(k\Delta t, l\Delta t) (s(t - k\Delta t)s(t - l\Delta t)) (\Delta t)^2 + \dots$$

$$\Rightarrow K_0 + \int K_1(\gamma) s(t - \gamma) d\gamma$$

$$+ \int K_2(\gamma_1, \gamma_2) s(t - \gamma_1) s(t - \gamma_2) d\gamma_1 d\gamma_2$$

$$- \underbrace{\Delta t P \int K_2(\gamma, \gamma) d\gamma}$$

Not what we expected. So we have to think of $s(t)$ as being a process definitely a constant P_{st} , i.e., std dev. of $s(t)$ is

$\sqrt{P_{st}}$. Then we'll get a well-defined limit;

i.e., if we keep $P = P_{st}$ constant. i.e., $P = \frac{1}{\Delta t} \int P$ on any realization.

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Calculating the a_k 's (+ the k 's).

$$\text{In general, } a_k = \frac{\langle f(x) \varphi_k(x) \rangle}{\langle \varphi_k(x) \varphi_k(x) \rangle} \quad [\text{Non-Hilbert space case}]$$

where $\varphi_k(x)$ is an orthogonal polynomial.

$$\text{For } a_k, \varphi_k(\vec{x}) = x_k; \quad a_k = \frac{\langle f(x) \cdot x_k \rangle}{\langle x_k^2 \rangle} = \frac{1}{P} \langle f(x) \cdot x_k \rangle$$

$$\text{In the continuum case, } \langle \varphi_k^2 \rangle = \Delta t \cdot P = P$$

$$\text{so } K(t) = \frac{\langle R(t) s(t-\tau) \rangle}{P}$$

$$\text{For } c_{kl}, \varphi_k(\vec{x}) = x_k x_l, \quad c_{kl} = \frac{\langle f(x) x_k x_l \rangle}{\langle x_k^2 x_l^2 \rangle} = \frac{1}{P^2} \langle f(x) x_k x_l \rangle$$

$$\langle \varphi_k^2 \rangle = (\Delta t P)^2 = P^2$$

$$K_2(t_1, t_2) = \frac{1}{Z P^2} \langle R(t_1) s(t-t_1) s(t-t_2) \rangle$$

\uparrow from $k \neq l \Rightarrow k < l$.

$$b_k: \varphi_k(\vec{x}) = x_k^2 - P, \quad \langle \varphi_k^2 \rangle = \langle x_k^4 - 2P x_k^2 + P^2 \rangle = Z P^2$$

$$b_k = \frac{1}{Z P^2} \langle f(x_k) (x_k^2 - P) \rangle$$

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$$K_2(\tau, \tau) = \frac{1}{2} p_2 \langle R(t) [s(t-\tau)^2 - p] \rangle$$

Unty $K_2(\tau, \tau)$ for $\tau_1 \neq \tau_2$ or $\tau_1 = \tau_2$.

The difference is when t_0 . Now $p = \frac{1}{dt} \int p$.

The subtracted term integrates to p (it is 0 except in a window of, which is $\frac{1}{dt} p$)

So we can write

$$K_2(\tau_1, \tau_2) = \frac{1}{2} p_2 \langle R(t) (s(t-\tau_1)s(t-\tau_2) - s(\tau_1-\tau_2)p) \rangle$$

Put another way, to calculate $K_2(\tau, \tau)$, we

cross correlate with $s(t_0-\tau) s(t-\tau) - s(\tau-\tau)$

for orthogonality.

To reconstruct the imp,

$$R(t) = K_0 + \int k_1(\tau) s(t-\tau) d\tau +$$

$$+ \int K_2(\tau, \tau_2) (s(t-\tau_1)s(t-\tau_2) - p s(\tau_1-\tau_2)) d\tau_1 d\tau_2$$

+ ...

The issues that arise in a practical implementation:

Choose ΔT . Choose $P = \frac{1}{\Delta T} P_0$.

Choose maximum time lag for analysis.

Estimate $\langle R(s) \cdot s(t-\tau) \rangle$ from finite samples of s .

$$\langle R(s) (s(t-\tau) s(t-\tau_0)) \rangle$$

Two: Main issue: samples of s may not be typical of the Gaussian noise, and therefore,

$$\text{the orthogonal factors } s(t-\tau), s(t-\tau_1) s(t-\tau_2) + s(t-\tau)^2 - P, \\ s(t-\tau_1)^2 - 3Ps(t-\tau_1), \text{ etc. may not}$$

be orthogonal w.r.t. the sample of noise.

So if, really, $R = \varphi_k'(s)$ but

if the estimate of $\langle \varphi_k(s) \varphi_l(s) \rangle \neq 0$
then

R will appear to have a component $\varphi_l(s)$.

So, can we choose "specific" samples to make the orthogonality as close as possible?

- Broaden the issue:
- (1) choose some other ensemble R of signals (not necessarily Gaussian white noise).
 - (2) constraint on orthogonal function series
 - (3) design inputs to sample R

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We could do fits with a sequence of R 's that approach Gaussian white noise -
 either analytically, or, from the system's point of view

or

we could even choose R based on a biological motivation - natural sounds,

natural sounds.

Relation to regression (functional, image, analysis)

We have $R(t)$ + several variables $v_1(t), \dots, v_n(t)$

and want to model $R(t)$ as a sum of "effects",

$$R(t) = \sum_{k=1}^n \sum_{\tau} \alpha_k(\tau) V_k(t - \tau).$$

$$\text{or even } \sum_{k=1}^n \sum_{\tau_1, \tau_2} \beta_k(\tau_1, \tau_2) [V_k(t - \tau_1) V_k(t - \tau_2)]$$

$$= \sum_{\tau} c_\tau V_\tau(t).$$

Want to orthogonalise the
 V_τ 's into ψ_τ 's.

Best if they were already orthogonal, bad if they are linearly dependent

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Two strategies for "a sequence of R's that approach GWN"

M-sequences "pseudorandom binary sequences"
 $s = \sum m_i \sin(\omega_i t + \phi_i)$

M-sequences approach GWN from the "system's point of view":

any real system has a front end that sums over-time,

$$\text{turning } s(t) \text{ into } s'(t) = \int f(\tau) s(t-\tau) d\tau$$

So even if $s(t)$ is only 0's & 1's, $s'(t)$ is \sim Gaussian.

Basic idea (but not M-sequences)

Say we have 3 time steps. 8 possible stimulus histories

0 0 0
0 0 1
0 1 0
0 1 1
1 0 0
(0 1
1 1 0
1 1 1

We could do 8 experiments, presenting each one several times (farsighted)

Or, $\overbrace{0 0 0 1 0 1 1 1 0 0 0 1 0 1 1 1 \dots}$
 present.

Note that past histories contain each of the 3-bit inputs;
 but we gain 3x in efficiency.

Basic idea is to choose a sequence of 0's & 1's for which the ϕ_i 's
 are nearly orthogonal.

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We will construct a sequence of length $2^n - 1$ which

(a) contains all n -tuples except $(0 \dots 0)$ exactly once

(b) is shift-orthogonal, i.e., $\langle \sigma(k), \sigma(k-n) \rangle = \pm \frac{1}{2^n - 1}$

(c) will allow for measurement of high-order kernels at short lags

(d) " " " " first-order kernels at all lags, which often is enough to characterize the system.

$\sigma: +1 \text{ or } -1$, it's helpful to think of it consisting of 0's & 1's

Recipe We will create a finite field of size 2^n , & write a table of logarithms w.r.t.

$n=4$.

Consider the finite field $\mathbb{F}_2 = \{0, 1\}$, and let's choose a polynomial of degree n that does not factor in \mathbb{F}_2 (other conditions too).

$$p(x) = x^4 + x + 1$$

This constructs a finite field, as follows:

Set $p(x)=0$, & multiply!

$$\begin{aligned} x^0 &= 1 \\ x^1 &= x \\ x^2 &= x^2 \\ x^3 &= x^3 \\ x^4 &= x+1 \\ x^5 &= x^2+x \\ x^6 &= x^3+x^2 \\ x^7 &= x^3+x+1 \quad \Rightarrow \text{E.g., } x^8 = x^7 \cdot x = (x^3+x+1) \cdot x \\ x^8 &= x^2+x+1 \\ x^9 &= x^3+x \\ x^{10} &= x^2+x+1 \\ x^{11} &= x^3+x^2+x \\ x^{12} &= x^3+x^2+x+1 \\ x^{13} &= x^3+x^2+1 \\ x^{14} &= x^3+1 \\ x^{15} &= 1 \end{aligned}$$

Then any column is the n -sequence.

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A recurrence rule holds within a column:

$$p(x) = x^n + g(x)$$

$$x^k \cdot x^n = x^k \cdot g(x)$$

$$x^{k+n} = \sum_{r=0}^{n-1} b_r x^{k+r}$$

$$g(x) = \sum_{r=0}^{n-1} b_r x^r$$

$$\text{Here, } b_0 = 1, b_1 = 1, b_2 = b_3 = 0$$

So, for example $k=5$ $b=0$ $b=1$ $b=2$ $b=3$ $b=4$	$x^5 =$ $x^6 =$ $x^7 =$ $x^8 =$ $x^9 =$	$x^2 + x$ $x^3 + x^2$ $x^3 + x + 1$ $x^2 + 1$ $x^3 + x$	    
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$$x^9 = x^6 + x^5$$

"Shift register" generation rule is the original table, equiv to this recursion.

If the sequence has its maximum length, then no n -tuples can repeat. (Otherwise it would close early).
And no 0-tuple.
So all n -tuples appear once.

What about shift-on-the-goability?

$$a \neq 0: \quad \underbrace{x^k}_{\text{const term}} + \underbrace{x^{k+a}}_{\text{const term}} = \begin{cases} 0 & \text{if match} \\ 1 & \text{if mismatch} \end{cases}$$

so $x^k + x^{k+a}$ is cross correlation of σ_k d σ_{k+a} .

$$x^k + x^{k+a} = x^k(1+x^a) = x^k \cdot x^{k+a}$$

because $(+x^a$ must be some $I(a)$) $x^k \cdot x^{k+a} = x^{k+I(a)}$.

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The above shows that on an n -sequence XOR'd a shift of itself is just another shift of the same n -sequence.

So $x^{k+2\alpha}$ has \sim equal # of 1's & 0's, so $x^k \cdot x^{k+\alpha}$ are independent

The drawback is that high-order regressors (polynomial) overlap with low-order ones:

$$\langle r(t) \sigma(t-\gamma) \sigma(t-\gamma_2) \rangle$$

$$\text{must} = \langle r(t) \sigma(t-\mathcal{I}(\gamma, \gamma_2)) \rangle$$

$$\text{since } \sigma(t-\gamma_1) \sigma(t-\gamma_2) = \sigma(t-\mathcal{I}(\gamma_1, \gamma_2))$$

by above argument

What can we do? Choose $p(x)$ so $\mathcal{I}(\gamma_1, \gamma_2)$ is large.

Or, let R contain $\{\sigma_k\}$ and $\{-\sigma_k\}$

so that now, $\sigma(t-\gamma)$ and $\sigma(t-\gamma) \sigma(t-\gamma_2)$

are orthogonal over R .

("Invert repeat" method)